

# Optimal Geospatial Query Placement in Cloud

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**Abstract.** Computing resources requirements are increasing with the massive generation of geospatial queries. These queries extract information from a large volume of spatial data. Placement of geospatial queries in virtual machines with minimum resource and energy wastage is a big challenge. Getting query results from mobile locations within a specific time duration is also a major concern. In this work, we have formulated a bi-objective optimization problem to minimize the energy consumption of cloud servers and service processing time. To solve the problem, we have proposed a crow search based bio-inspired heuristic. We compared our proposed algorithm through simulation with traditional First Fit, and Best Fit algorithms and the obtained results are significantly better than the traditional techniques.

**Keywords:** Geospatial Query, Cloud Computing, Energy Efficiency

## 1 Introduction

Geospatial Query (GQ) processing is essential in the applications of geographical information systems (GIS), multimedia information systems (MIS), location-based services (LBS), etc. In GQ, the *location* is an essential attribute. To access GIS data, a user can generate GQ from their mobile devices either from a static location or mobile location. In the best case, if the query resolved in a nearby cloud data centre (DC), which leads to less communication cost and propagation delay. These metrics will not change for GQs, which are accessed from a static location. Whereas GQs are coming from mobile location users, metrics will change very frequently. GQs extract information from a large volume of spatial data [1]. Therefore, the distribution of GQs in dynamic cloud DCs is challenging [2]. Cloud computing offers a shared pool of huge resources like CPU core, RAM, storage, bandwidth, etc. Virtualization made physical resources available to the user by offering an illusion of a dedicated system. This technique creates multiple virtual machines (VMs) with different configurations of CPU cores, RAM, storage, bandwidth, etc. As per users requirement, the cloud service provider (CSP) offers them required VMs.

GQ load distribution means shifting of GQs to the heavy load VMs to low load VMs. Make a stable cloud environment with equal GQ distribution. While shifting the GQ, keep in mind that the query should be resolved within its timespan and the energy consumption will be minimum. In the literature, very fewer numbers of works published in GQ placement in cloud. Lee et al. [3] proposed a spatial indexing technique that works over HBase for big geospatial data. They considered containedIn, intersects, withinDistance types of GQ and used GeoLife and SFTaxi Trajectories data points for experiments. In [4], Bai et al. proposed an indexing technique over HBase distributed database using kNN and window query processing algorithms to process huge data objects. A learning technique to resolve GQ in cloud is discussed in [5]. Akdogan et al. [6] proposed Voronoi diagram for the efficient processing of a varied range of GQs. The authors have deployed MapReduce-based approach to resolve reverse nearest neighbor (RNN), maximum reverse nearest neighbor (MaxRNN) and k-nearest neighbor (kNN) queries. GQ resolution on the cloud using geospatial service chains has been discussed in [7].

On the other hand, many algorithms are proposed for load balancing in cloud. Load balancing is a key factor for GQ placement in cloud servers. Kumar and Sharma [8] proposed an algorithm that is based on a proactive prediction based approach. It predicts future loads based on past load history and distributed the loads from heavy load VMs to underloaded VMs. Calheiros et al. [9] proposed a workload prediction using the ARIMA model, which helps to get feedback from previous workloads and update the model accordingly. It helps to assign VMs dynamically and maintain the user QoS by reducing the rejection rate and response time. Garg et al. [10] proposed a mechanism to dynamic cloud resource allocation by assigning maximum workload to a DC as per SLA. To maximize the utilization of the cloud DC, they integrated non-interactive and transactional applications. It also reduces the penalty due to less SLA violation as a maximum number of the applications are run in DCs. [11] proposed a framework which schedules tasks in an energy-efficient way. It decreases the overall energy of cloud data center by predicting similar kind of workload distribution as actual one.

From the above literature survey, we observe that most of the works are done separately for GQ processing and load distribution in the cloud environment. We have merged these two scenarios. In this work, GQs are not only resolved in the user-specified timespan but also minimize the overall power consumption of the cloud environment due to an efficient GQ load balancing algorithm among the VMs. In this paper, we have assumed that the scheduling of GQs is highly heterogeneous. We have focused on the minimization of GQ resolving timespan and energy consumption. The key contributions of this paper are:

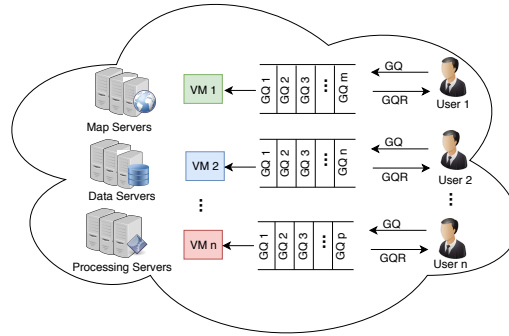
- Optimal GQ load distribution with minimal timespan and energy consumption.
- Minimize the service time of the GQ.

The rest of the paper is organized as follows. Categories of the geospatial queries and modeling of processing those queries in the cloud platform are explained in

Section 2. In Section 3, we have discussed the problem-solution approach using crow search heuristic. Performance analysis of the proposed scheme is presented in Section 4. Conclusion and future scope of our work are drawn in the last section.

## 2 System Model

Users generate the GQs through the user interface of web-enabled electronic gadgets, i.e. mobile, laptop, computer, etc. It will be submitted to the Cloud broker. The cloud broker will map the GQ with the existing query types. It is also needed to identify the requirements of geospatial data and geospatial services. There are three types of servers i.e. Processing Server, Data Server, and Map Server available in the DC. The data server keeps geospatial data. The processing server processes the geospatial data. The map server helps to generate the map on processed geospatial data. After identification, the broker assigns a VM to execute GQ. When the execution is over, it generates the GQ results (GQR), which projected to the user interface. A pictorial view of our system model has shown in Fig. 1.



**Fig. 1.** Geospatial Query Processing Model in Cloud

Assignment of appropriate VM to a GQ is a key feature to distribute the GQ load. Before the assignment of GQ, it is needed to know GQ types, which helps to select VM's specifications during VM to server mapping. The types of GQs [1] are mentioned below.

- **Filter Query**- This type of query [12] filters a particular geometry which presents in the another geometry.
- **Within Distance**- It measures whether one geometry or object is present within a particular euclidean distance of another geometry or not.

- **Nearest Neighbour (NN)**- It measures whether geometries is the nearest neighbor of a particular geometry or not.
- **Geospatial Join Query**- It compares one layer of a geometry with the layers of the other geometries. Geospatial index type (that is, R-tree or Quadtree) must be the same on the geometry column of all the tables involved in the join operation.

## 2.1 Service configuration

Cloud DC receives GQs from end-users. To assign GQs in VMs, it has to meet the resource requirements such as CPU, memory, etc. The VM assigns to the GQ iff it meets the resource requirements of GQ. The selection of VMs with a number of GQs will be based on its capacity constraints. If requested VM specification is not available, then CSP may assign a higher configured VM. A cloud DC is configured with a large number of servers. VMs are also assigned to the nearby cloud DC of GQ. The nearby cloud DC should fulfill the resource requirement of GQ. Therefore, the optimal mapping of VM to the cloud server is important for maintaining the QoS of GQ services. Scheduling of the GQs is required to place into VMs. After the completion of GQ scheduling, an algorithm is needed for optimal selection of VM among all available VMs.

## 2.2 VM to cloud server mapping

A suitable mapping of GQs to VMs is required while optimizing overall GQ processing time and power consumption of the system. A geospatial query set  $GQ = \{GQ_1, GQ_2, \dots, GQ_p\}$  consists of ‘ $p$ ’ number of geospatial queries. A VM set  $\{vm_1, vm_2, \dots, vm_q\}$  consists of ‘ $q$ ’ VMs. Each GQ has two dimensional (CPU and memory) resource requests. We have generated a crow matrix for the GQ requests set. The main focus of our GQ placement (GQP) algorithm is to the continuous optimization of processing time and power consumption. Next, we have represented two parameters mathematically.

**Processing time calculation:** The service time of the GQ processing can be defined as follows:

$$T_s = T_w + T_p \quad (1)$$

Where,  $T_s$  is Service Time,  $T_w$  is Waiting Time in Queue, and  $T_p$  is Processing Time in VM. We consider the queue discipline as m/m/1 [13]. All GQ requests will be in a FIFO manner. The arrival rate of GQs is  $\lambda$  and  $\mu$  will be the service rate. We can calculate the steady-state probability of ‘ $p$ ’ numbers of GQs. Thus, we can calculate the expected number GQ request in the queue, by which the  $T_w$  in the queue can also be calculated. Similarly, we can calculate the  $T_s$ .

Now, to process large number of GQs, it is needed more VMs. If we increase the number of VMs, then the energy will consume more. This is an NP-Hard problem. We have to trade-off between the number of VMs generation and overall energy consumption. The relation between processing time in VM ( $T_p$ ) and

energy consumption

$$T_p \propto E_c \quad (2)$$

**Power consumption calculation:** The total power consumption can be defined as below [13]:

$$pwr_i = (pwr_i^{max} - pwr_i^{min}) * utz_i + pwr_i^{min} \quad (3)$$

$pwr_i^{max}$  and  $pwr_i^{min}$  are the maximum and minimum average power consumed while maximum and minimum utilization is occurred respectively.  $utz_i$  is utilization of  $vm_i$ .

### 2.3 Problem formulation

In this paper, the GQ placement algorithm is modelled as a bi-objective optimization problem. Our objective is to minimize the energy consumption of cloud servers and processing time of GQ. The objective function of the GQP algorithm is mentioned below:

$$\text{Minimize } \sum_{c=1}^n E_c \quad \text{and} \quad \text{Minimize } \sum_{p=1}^m T_p \quad (4)$$

Subject to the three constraints are mentioned below:

- **Assignment constraint:** It assures that the GQs are placed in such VMs where each GQ dimensions are matched with the VM dimensions.
- **Capacity constraint:** It assures that the total VM requirements of the GQ set should be less than or equal to the total available VMs.
- **Placement constraint:** It assures to the assignment of a GQ to only one VM which meets the resource requirements along with all dimensions.

## 3 Crow Search Algorithm for Geospatial Query Placement

To solve the aforementioned bi-objective optimization problem, we choose the crow search algorithm (CSA) [14]. CSA is a well-known algorithm that is used to solve many optimization problems. Unlike genetic algorithm (GA), ant colony optimization (ACO), particle swarm optimization (PSO), chaotic ant swarm (CAS), CSA also makes use of a population of seekers to explore the search space. In CSA, the number of adjustable parameters is less (flight length (FL), and awareness probability (AP)) compare to the other optimization algorithms. As adjustable parameters are very difficult to manage. GQP is such an optimization problem where we can use CSA to find its optimal solutions. For GQ placement, GQs are considered as crows, and the best VM for GQ placement is equivalent to an optimal food source. A suitable VM selection is needed as all VMs are uniformly eligible for placing the GQs. As crows search for optimal food sources, similarly, the GQs are searching for appropriate VM for processing. Our proposed crow search based algorithm for geospatial query placement into VM is described in Algorithm 1.

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**Algorithm 1:** Crow search algorithm for GQ placement (GQP).

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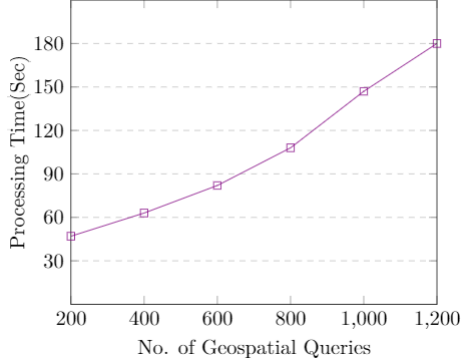
**Input:** GQs from different users  
**Result:** Assign GQs in appropriate VM  
Evaluate the position of the GQs.  
Initialize the memory of the GQs.  
**while**  $iter < iter_{max}$  **do**  
    Generation of a new assignment of GQs in set  $GQS''$ .  
    **for**  $i = 1$  to  $n$  **do**  
        Randomly select a VM  $vm_j$  that  $gq_i$  follows.  
        Define the awareness probability AP.  
        **if**  $r_j \geq AP_j^{iter}$  **then**  
            Check the availability of the resources at VM  $vm_j$ .  
            Update the position of  $gq_i$  in an array  $GQS'$ .  
        **else**  
            Generate a random VM  $vm_k$  with resources.  
            Update the position of  $gq_i$  in the array  $GQS'$ .  
        **end**  
    **end**  
    **if**  $FL < 1$  **then**  
        Set  $GQS'' = GQS'$ .  
    **else**  
        Perform 'x' random shuffles  
        Update  $GQS''$ .  
    **end**  
    **if**  $f(GQS'')$  is better than  $f(GQS)$  **then**  
        Set  $GQS = GQS''$ .  
    **end**  
**end**

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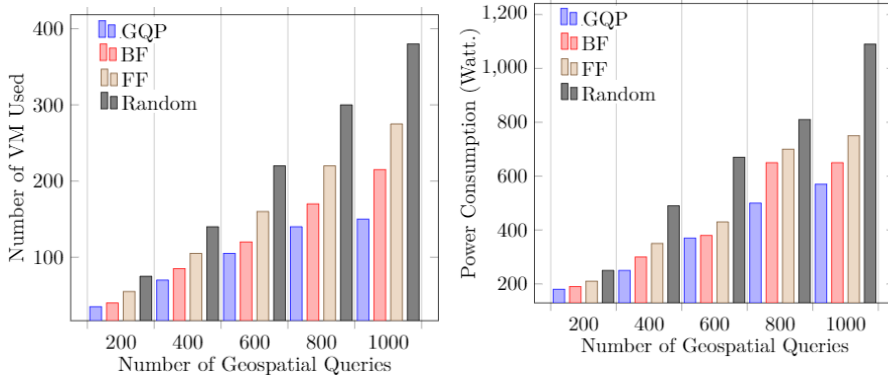
## 4 Performance Evaluation

We have performed experiments in CloudSim 4.0 simulator [15]. The considered number of hosts 100 and each host has 16 processing elements. Number of VMs is varying from 100 to 500, and five types of VMs are considered. These are *Micro VM* (1 Core, 1 GB RAM), *Small VM* (1 Core, 2 GB RAM), *Medium VM* (2 Core, 2 GB RAM), *Large VM* (2 Core, 8 GB RAM), and *xLarge VM* (1 Core, 1 GB RAM). We have taken the number of cloudlets (here GQs) within the range of 200-1200. For keeping the complexity of the algorithm less, we considered the value of random shuffle is 20, and the value of  $iter_{max}$  is 1000.

We have processed the GQs in the first-come, first-serve basis. The number of GQs against processing time graph has been displayed in Fig. 2. Also, we have compared our GQP algorithm with existing Best Fit (BF), First Fit (FF) and Random allocation algorithms in the context of the number of used VM with respect to the number of geospatial queries which are depicted in Fig. 3. It has been observed that the number of used VMs are increased with the increment of the number of GQs. The number of VMs are used in GQP is lesser than the other three existing algorithms. In FF, first GQ is placed in the first VM. Next GQ will check whether it will fit in the first VM or not. If not, then it moves to the later VM. This moves toward non-optimal solutions. In the case of BF, the algorithm checks all the VMs capacity and then decides where to move the GQs. This is also moved toward non-optimal solutions. Fig. 4 shows the overall power consumption against the different types of GQs placement strategies. As less number of VMs are used for resolving GQs in our GQP algorithm, this



**Fig. 2.** Processing Time vs Number of Geospatial Queries



**Fig. 3.** VMs used by number of GQs for different placement strategies

**Fig. 4.** Variation of power consumption with the number of GQs for different placement strategies

leads to the minimal power consumption in GQP compared to the other three algorithms.

## 5 Conclusions and Future Work

The processing of GQ from a mobile location in the cloud server is a challenging task. Mainly, GIS data resides in the spatial database in huge volumes. To process a large number of GQ leads to high propagation delay and response time. Also, it causes higher energy consumption in cloud servers. In this work, we have formulated an optimization problem and solved it using crow search based heuristic. The obtained results are significantly outperforming than traditional FF and BF. As part of an immediate extension of this work, we will implement GQ processing in a multi-cloud environment such as cloud federation. We will

verify how the performance will change from a single cloud to a multi-cloud environment.

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